

Problem Set One

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1. For each claim, determine whether the statement is **True** or **False**. Justify your answer.

- (a) $n + 3 = O(n^3) \rightarrow$ **True**. According to the definition of Big-O notation, $f = O(g)$ if

$$(\exists c, k > 0, x > k) [|f(x)| \leq c |g(x)|]$$

Therefore

$$|n + 3| \leq c |n^3|$$

and the statement is valid.

- (b) $3^{2n} = O(3^n) \rightarrow$ **False**. Again, using the previous definition of Big-O notation we see that

$$|3^{2n}| \not\leq c |3^n|$$

- (c) $n^n = o(n!) \rightarrow$ **False**. We can use the definition of little-o notation which states that f is little-o of g if

$$\lim_{x \rightarrow \infty} \frac{f(x)}{g(x)} = 0$$

Therefore this statement turns to

$$\lim_{x \rightarrow \infty} \frac{n^n}{n!} = \infty$$

Therefore the statement is invalid.

- (d) $\frac{1}{3^n} = o(1) \rightarrow$ **True**. We then use the above definition again and apply L'Hospital's Rule to determine the value of the limit.

$$\lim_{x \rightarrow \infty} \frac{1}{3^n} \rightarrow \lim_{x \rightarrow \infty} \frac{0}{3} = 0$$

- (e) $\ln^3(n) = \Theta(\log_2^3(n)) \rightarrow$ **False**¹ We can use the definition of Big-O notation to determine that

$$\ln^3(n) = O(\log_2^3(n))$$

¹This is assuming that $lg(x)$ refers to the base-2 logarithm, $\log_2(x)$.

because $|\ln^3(n)| \leq c |\log_2^3(n)|$, however

$$\boxed{\log_2^3(n) \neq O(\ln^3(n))}$$

because $|\log_2^3(n)| \not\leq c |\ln^3(n)|$. Therefore the statement is false.

2. Simplify each of the following expressions.

(a)

$$\frac{d}{dt}(3t^4 + 1/3t^3 - 7) \rightarrow \boxed{12t^3 + t^2}$$

(b)

$$\sum_{i=0}^k 2^i \rightarrow 1 + 2 + 4 + \dots + 2^k \rightarrow \boxed{2^{k+1} - 1}$$

(c)

$$\Theta\left(\sum_{k=1}^n \frac{1}{k}\right) \rightarrow \boxed{H_n}$$

Where H_n is the n^{th} Harmonic number.

3. T is a balanced binary search tree storing n values. Describe an $O(n)$ -time algorithm that takes input T and returns an array containing the same values in ascending order.

(a) Below is the code to perform this operation.

```

1  asc = [] # List to populate
2  class Node: # The structure of any given node
3      left = None # Class object of left node
4      right = None # Class object of right node
5      value = None # Value of node
6  def tree_to_array(head): # Function to scrape in asc order
7      if head.left != None: # If left is node
8          tree_to_array(head.left) # Take left
9          head.left = None # Destroy traversed result
10     if head.right != None: # Else take right
11         asc.append(head.value) # Take next smallest val
12         tree_to_array(head.right) # Go right
13         head.right = None # Destroy traversed result
14     if head.left is None and # If both sides are empty
15         head.right is None:
16         try:
17             if head.value >=
18                 asc[len(asc) - 1]: # If larger than prev
19                 asc.append(head.value) # This value is our next smallest
20         except IndexError: # Only enter if list is empty
21             asc.append(head.value) # This value is our next smallest
22 head = construct_tree(random=true) # Create a random balanced tree
23 print(tree_to_array(head)) # Print our end array

```

4. Acme Corp. has asked Professor Flitwick to develop a faster algorithm for their core business. The current algorithm runs in $f(n)$ time. (For concreteness, assume it takes $f(n)$ microseconds to solve a problem of size exactly n .) Flitwick believes he can develop a faster algorithm, which takes only $g(n)$ time, but developing it will take t days. Acme only needs to solve a problem of size n once. Should Acme pay Flitwick to develop the faster algorithm or should they stick with their current algorithm? Explain.

(a) Let $n = 41$, $f(n) = 1.99^n$, $g(n) = n^3$ and $t = 17$ days.

i. The time it will take the original algorithm to complete is

$$1.99^n \text{ where } n = 41 \rightarrow 1790507451731.9128ms \rightarrow 20.7235d$$

Flitwick can complete and run his algorithm in

$$17 + n^3 \text{ where } n = 41 \rightarrow 17d + 68921ms \rightarrow 17.0000007977d$$

Therefore the company *should* pay him to develop the better algorithm as it will save them 3 days time.

(b) Let $n = 10^6$, $f(n) = n^{2.00}$, $g(n) = n^{1.99}$ and $t = 2$ days.

i. The time it will take the original algorithm to complete is

$$n^{2.00} \text{ where } n = 10^6 \rightarrow 1000000000000ms \rightarrow 11.5741d$$

Flitwick can complete and run his algorithm in

$$2 + n^{1.99} \text{ where } n = 10^6 \rightarrow 2d + 870963589956.0806ms \rightarrow 12.0806d$$

Therefore the company *should not* pay him to develop the better algorithm as it will take an extra day and a half to complete.

5. Using the mathematical definition of Big-O, answer the following. Show your work.

(a) Is $2^{nk} = O(2^n)$ for $k > 1$?

i. No. 2^{nk} will always grow faster than 2^n .

$$2^{nk} \rightarrow (2^n)^k \rightarrow |(2^n)^k| \not\leq c|2^n|$$

(b) Is $2^{n+k} = O(2^n)$, for $k = O(1)$?

i. Yes. 2^k is constant, therefore

$$2^{n+k} \rightarrow 2^n 2^k \rightarrow |2^n 2^k| \leq c|2^n|$$

6. Is an array that is in sorted order also a min-heap? Justify.

(a) Technically no, they are not the same. They have differing data structures, however they are more similar than not upon further inspection. A sorted array has the form [1, 2, 3, 4, 5] while a min-heap has the form

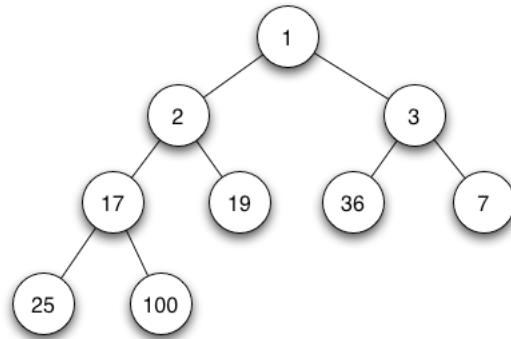


Figure 1: A Sample Min-Heap

with a corresponding data structure similar to the sample code below.

```

1 class Node:
2     left = left_node_class_object # Must be greater than Node
3     right = right_node_class_object # Must be greater than Node
4     value = node_value
    
```

As is evident the fundamental data structures expressing the two are not similar in the slightest. This being said however, a sorted array will correspond to the following min-heap

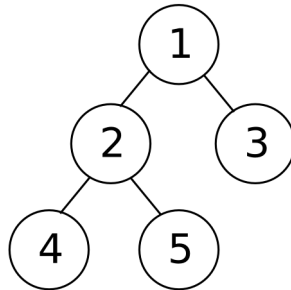


Figure 2: The Min-Heap for our Array

When the min-heap is accessed top-down, left-to-right it will have a one-to-one correspondence to our array. So to put it succinctly, the two data structures are not the same, however they have similar appearance and behavior.